

Instruction latencies and throughput for AMD and Intel x86 processors

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This report is work-in-progress. A newer version might be available here: <https://gmplib.org/~tege/x86-timing.pdf>

In this short report we present latency and throughput data for various x86 processors. We only present data on integer operations. The data on integer MMX and SSE2 instructions is currently limited. We might present more complete data in the future, if there is enough interest.

There are several reasons for presenting this report:

1. Intel's published data were in the past incomplete and full of errors.
2. Intel did not publish any data for 64-bit operations.
3. To allow straightforward comparison of an important aspect of AMD and Intel pipelines.

The here presented data is the result of extensive timing tests. While we have made an effort to make sure the data is accurate, the reader is cautioned that some errors might have crept in.

1 Nomenclature and notation

LNN means *latency for NN-bit operation*. *TNN* means *throughput for NN-bit operation*.

The term *throughput* is used to mean *number of instructions per cycle of this type that can be sustained*. That implies that more throughput is better, which is consistent with how most people understand the term. Intel use that same term in the exact opposite meaning in their manuals.

The notation "P6 0-E", "P4 F0", etc, are used to save table header space. P4 means Pentium 4, which has family number F (in hexadecimal). Then there are model numbers 0, 1, 2, etc, where 4 and above means the core has the AMD defined 64-bit instructions. P6 means family 6, which is a big happy family with Pentium Pro, Pentium 2, Pentium 3, Pentium M, Core and Core 2, and any other Intel processor from 2010 onwards.

In some of these tables, Core 2 numbers hide under the P6 F moniker (family 6, model F).

Family and model numbers are returned by the `cpuid` instruction.

2 How data was generated

The throughput numbers here have been generated using a loop with the measured instruction repeated many times. This results in some cases in lower numbers than those claimed by the processor vendors. Our measure-

ment method requires that the entire pipeline (cache, decode, issue, execution, complete) can sustain the indicated execution rate.

Measuring performance for immediate operands is tricky, as x86 encodes small immediate operands specially, and as different cores are optimised for different ranges.

An example of how complex this can be is that Pentium F0-F2 can sustain 3 `add r,i` per cycle for $-32768 \leq i \leq 32767$, but for larger immediate operands it can sustain only about 3/2 per cycle. Pentium F3 behaves similarly, but for $-65536 \leq i \leq 65535$. This is not related to x86 instruction encoding; the exact same encoding is used for greater immediate operands. (Operands that fit into a byte have a special encoding, though.)

3 Comments on table data (this section is mostly obsolete)

The Pentium 4 performance for 64-bit right shifts is really poor. 64-bit left shift as well as all 32-bit shift have acceptable performance. It appears that the data path from the upper 32 bits to the lower 32 bit of the ALU, is not well designed.

On Pentium 4, it is possible to reach much better 32-bit `shl/sal/shr/sar r,c` performance if dummy shift instructions with immediate counts are strategically inserted. The best measured resulting performance is about 1.4 instructions/cycle.

On AMD K8, the odd man out is `test r,i`. The reason is that the encoding of this instruction doesn't provide any short immediate version. The L1 instruction cache can provide only 16 bytes per clock, but 3 `test r,i` need 18 or 21 bytes (depending on register). The AMD K10 has a wider bus between the L1 instruction cache and decoders, and can sustain 3 insns/cycle even for these long instructions. (One might think it is more than a little silly to measure three consecutive `test` instructions, since these instructions are pointless without a branch or set-on-condition instruction. But it gives some idea of the resource usage of these instructions, and we prefer to keep the measurements orthogonal between instructions.)

The data for `adc` and `sbb` are approximate. A long chain of these instructions would have a recurrent dependency on the carry flag. Our sequences break such dependencies by regularly executing a plain `add`, relying on out-of-order execution and carry flag renaming. It is likely that slightly better performance could be achieved than what is here indicated.

Sequences with `adc` and `sbb` interleaved with `inc` or `dec` run poorly on all P6 processors (including Core processors), they result in a pipeline hiccup of about 12 cycles. That sort of problems cannot be represented in the table numbers.

4 Your feedback

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		Intel Atom		Intel SLM		Intel GLM		Intel GLM+		AMD BT1		AMD BT2		VIA Nano	
		L64	T64	L64	T64	L64	T64	L64	T64	L64	T64	L64	T64	L64	T64
add	r,ri	1	2	1	2	1	2	1	3	1	2	1	2	1	2
sub	r,ri	1	2	1	2	1	2	1	3	1	2	1	2	1	2
and	r,r	1	2	1	2	1	2	1	3	1	2	1	2	1	2
or	r,r	1	2	1	2	1	2	1	3	1	2	1	2	1	2
xor	r,r	1	2	1	2	1	2	1	3	1	2	1	2	1	2
inc	r	1	2	1	1	1	1	1	1	1	2	1	2	1	2
dec	r	1	2	1	1	1	1	1	1	1	2	1	2	1	2
neg	r	1	2	1	2	1	2	1	3	1	2	1	2	1	2
not	r	1	2	1	2	1	2	1	3	1	2	1	2	1	2
imul	r,ri	14	1/11	5	1/2	5	1/2	5	1/2	6	1/4	6	1/4	5	1/2
mul	r	18	1/18	7	1/8	5	1/2	5	1/2	6-7	1/5	6-7	1/5	12	1/2
mulx	r,r,r	-	-	-	-	-	-	-	-	-	-	-	-	-	-
div	r	191		95		42		42		81		44		157	
adc	r,ri	2	1/2	2	1/2	2	1/2	2	1/2	1	1	1	1	1	1
sbb	r,ri	2	1/2	2	1/2	2	1/2	2	1/2	1	1	1	1	1	1
ad*x	r,r,r	-	-	-	-	-	-	-	-	-	-	-	-	-	-
shl	r,i	1	1	1	1	1	1	1	1	1	2	1	2	1	1
sal	r,i	1	1	1	1	1	1	1	1	1	2	1	2	1	1
shr	r,i	1	1	1	1	1	1	1	1	1	2	1	2	1	1
sar	r,i	1	1	1	1	1	1	1	1	1	2	1	2	1	1
shl	r,c	1	1	1	1	2	1	1	1	1	2	1	2	1	1
sal	r,c	1	1	1	1	2	1	1	1	1	2	1	2	1	1
shr	r,c	1	1	1	1	2	1	1	1	1	2	1	2	1	1
sar	r,c	1	1	1	1	2	1	1	1	1	2	1	2	1	1
shld	r,r,i	11	1/11	10	1/10	12	1/12	11	1/11	4	1/3	4	1/3	35	1/35
shrd	r,r,i	9	1/9	10	1/10	12	1/12	11	1/11	4	1/3	4	1/3	45	1/45
shld	r,r,c	10	1/10	10	1/10	14	1/14	13	1/13	4	1/4	4	1/4	35	1/35
shrd	r,r,c	9	1/9	10	1/10	14	1/14	13	1/13	4	1/4	4	1/4	45	1/45
shlx	r,r,r	-	-	-	-	-	-	-	-	-	-	-	-	-	-
shrx	r,r,r	-	-	-	-	-	-	-	-	-	-	-	-	-	-
rol	r,i	1	1	1	1	1	1	1	1	1	2	1	2	1	1
ror	r,i	1	1	1	1	1	1	1	1	1	2	1	2	1	1
rol	r,c	1	1	1	1	2	1	1	1	1	2	1	2	1	1
ror	r,c	1	1	1	1	2	1	1	1	1	2	1	2	1	1
cmp	r,ri	1	2	1	2	1	2	1	3	1	2	1	2	1	2
test	r,i	1	2	1	2	1	2	1	2	1	1.75	1	1.75	1	2
test	r,r	1	2	1	2	1	2	1	3	1	2	1	2	1	2
bt	r,i	2	1	1	1	1	1	1	1	1	2	1	2	2	1/2
mov	r,r	1	2	1	2	1	2	1	3	1	2	1	2	1	1
cmov	r,r	2	1/2	2	1	2	1	2	1	1	2	1	2	2	1
movzx	r,r	1	1	1	1	1	1	1	1	1	2	1	2	1	1
movsx	r,r	1	1	1	1	1	1	1	1	1	2	1	2	1	1
movsxd	r,r	1	1	1	1	1	1	1	1	1	2	1	2	1	1
bswap	r	1	1	1	1	1	1	1	1	1	2	1	2	1	1
lea	r,r+r	1	1	1	1	1	1	1	1	1	2	1	2	1	1
lea	r,b+r			1	2	1	2	1	3	1	2	1	2		
lea	r,r+r*s	1	1	1	1	1	1	1	1	2	1	2	1	1	1
lea	r,b+r+r	1	1	2	1	2	1	2	1	2	1	2	1	1	1
lea	r,b+r+r*s	1	1	2	1	2	1	2	1	2	1	2	1	1	1
bsr	r,r	17	1/16	10	1/10	10	1/8	9	1/8	8	1/8	4	1/4	3	1/2
bsf	r,r	16	1/16	10	1/10	10	1/8	9	1/8	11	1/8	4	1/4	3	1/2
lzcnt	r,r	-	-	-	-	-	-	-	-	6	1/6	1	2	-	-
tzcnt	r,r	-	-	-	-	-	-	-	-	-	-	2	1	-	-
popcnt	r,r	-	-	3	1	3	1	3	1	6	1/5	1	2	-	-
		Intel Atom		Intel SLM		Intel GLM		Intel GLM+		AMD BT1		AMD BT2		VIA Nano	

Table 2. Integer register instructions, 64-bit operations, low-end CPUs.

	Intel P4 F4		Intel Core 2		Intel NHM		Intel SBR		Intel HWL		Intel BWL		AMD K8-K9		AMD K10		AMD BD2		Intel Atom		VIA Nano	
	L32	T32	L32	T32	L32	T32	L32	T32	L32	T32	L32	T32	L32	T32	L32	T32	L32	T32	L32	T32	L32	T32
add r,ri	1	2.5	1	3	1	3	1	3	1	4	1	3.5	1	3	1	3			1	2	1	2
sub r,ri	1	2.5	1	3	1	3	1	3	1	4	1	3.5	1	3	1	3			1	2	1	2
and r,r	1	2	1	3	1	3	1	3	1	4	1	3.5	1	3	1	3			1	2	1	2
or r,r	1	2	1	3	1	3	1	3	1	4	1	3.5	1	3	1	3			1	2	1	2
xor r,r	1	2	1	3	1	3	1	3	1	4	1	3.5	1	3	1	3			1	2	1	2
inc r	1	1	1	3	1	3	1	3	1	4	1	3	1	3	1	3			1	2	1	2
dec r	1	1	1	3	1	3	1	3	1	4	1	3	1	3	1	3			1	2	1	2
neg r	1	2	1	3	1	3	1	3	1	4	1	3.5	1	3	1	3			1	2	1	2
not r	1	1.7	1	3	1	3	1	3	1	4	1	3.5	1	3	1	3			1	2	1	2
imul r,ri	10	1	3	1	3	1	3	1	3	1	3	1	3	1	3	1			5	1/2	4	1
mul r	11	1/2	5	2/3	1/2	4	1/2	4	1/2	4	1/2	4	1/2	3	1	3	1		9	1/9	8	1/2
mulx r,r,r	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-
div r	80 ²	1/34	40		26	26	28	31	39	1/39	45	1/45						50		39		
adc r,ri	10	1/3	2	1	2	1	2 ⁷	1	2 ⁷	1	1	1	1	3	1	3			2	1/2	1	1
sbb r,ri	10	1/3	2	1	2	1	2 ⁷	1	2 ⁷	1	1	1	1	3	1	3			2	1/2	1	1
shl r,i	1	1.7	1	2	1	2	1	2	1	2	1	2	1	3	1	3			1	1	1	1
sal r,i	1	1.7	1	2	1	2	1	2	1	2	1	2	1	3	1	3			1	1	1	1
shr r,i	1	1.7	1	2	1	2	1	2	1	2	1	2	1	3	1	3			1	1	1	1
sar r,i	1	1.7	1	2	1	2	1	2	1	2	1	2	1	3	1	3			1	1	1	1
shl r,c	2	1/2	1	2	1	2	2	1/2	2	1/2	2	1/2	2	1/2	1	3	1	3	1	1	1	1
sal r,c	2	1/2	1	2	1	2	2	1/2	2	1/2	2	1/2	2	1/2	1	3	1	3	1	1	1	1
shr r,c	2	1/2	1	2	1	2	2	1/2	2	1/2	2	1/2	2	1/2	1	3	1	3	1	1	1	1
sar r,c	2	1/2	1	2	1	2	2	1/2	2	1/2	2	1/2	2	1/2	1	3	1	3	1	1	1	1
shld r,r,i	8	1/7	2	1	4	1	1	2	1	1	1	1	3	1/2	3	1/2			5	1/5	7	1/7
shrd r,r,i	8	1/7	2	1	4	1	1	2	1	1	1	1	3	1/2	3	1/2			5	1/5	7	1/7
shld r,r,c	9	1/8	2	1	4	1	2	2/3	4	1/2	4	1/2	3	1/3	3	1/3			5	1/5	7	1/7
shrd r,r,c	9	1/8	2	1	4	1	2	2/3	4	1/2	4	1/2	3	1/3	3	1/3			5	1/5	7	1/7
shlx r,r,c	-	-	-	-	-	-	-	1	2	1	2	-	-	-	-	-	-	-	-	-	-	-
shrx r,r,c	-	-	-	-	-	-	-	1	2	1	2	-	-	-	-	-	-	-	-	-	-	-
rol r,i	1	1	1	1	1	1	1	1	1	2	1	2	1	3	1	3			1	1	1	1
ror r,i	1	1	1	1	1	1	1	1	1	2	1	2	1	3	1	3			1	1	1	1
rol r,c	2	1/2	1	1	1	1	2	1/2	2	1/2	2	1/2	1	3	1	3			1	1	1	1
ror r,c	2	1/2	1	1	1	1	2	1/2	2	1/2	2	1/2	1	3	1	3			1	1	1	1
cmp r,ri	1	2.5					3	4	4	4	1	3	1	3	1	3			1	2	1/2	2
test r,i	1	1.7					3	4	4	4	1	3	1	3	1	3			1	2	1/2	2
test r,r	1	2					3	4	4	4	1	3	1	3	1	3			1	2	1/2	2
bt r,i	8	1/8	2	1	2	1	2	1	2	1	1	1	1	3	1	3			2	1	2	1/2
mov r,r	1	2.5	1	3	1	3	1	3	1	4	1	4	1	3	1	3			1	2	1	1
cmov r,r	10	1	2	1/2	2	1	2	1	2	1	1	1	1	3	1	3			2	1/2	2	1
movzx r,r	1	2.5	1	3	1	3	1	3	1	4	1	4	1	3	1	3			1	1	1	1
movsx r,r	2	1.5	1	3	1	3	1	3	1	4	1	4	1	3	1	3			1	1	1	1
movsxd r,r	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-
bswap r	1	2	4	1	1	1	1	1	1	2	1	2	1	3	1	3			1	1	1	1
lea r,r+r			1	1	1	1							1	3	1	3			1	1	1	1
lea r,r+r*s			1	1	1	1							2	3	2	3			1	1	1	1
lea r,b+r+r			1	1	1	1							2	3	2	3			1	1	1	1
lea r,b+r+r*s			1	1	1	1							2	3	2	3			1	1	1	1
bsr r,r	16	1/2	2	1	3	1	3	1	3	1	3	1	11	1/10	4	1/3			17	1/16	3	1/2
bsf r,r	16	1/2	2	1	3	1	3	1	3	1	3	1	8	1/8	4	1/3			16	1/16	3	1/2
lzcnt r,r	-	-	-	-	-	-	-	3	1	3	1	-	-	2	1			-	-	-	-	-
popcnt r,r	-	-	-	-	3	1	3	1	3	1	3	1	-	-	2	1			-	-	-	-

Table 3. Integer register instructions, 32-bit operations on 64-bit CPUs.

	Intel P4						Intel P6		AMD		
	F0-F1		F2		F3-F4		O-E		K7		
	L32	T32	L32	T32	L32	T32	L32	T32	L32	T32	
add	r,ri	1/2	3	1/2	3	1	2.5	1	2	1	2.7
sub	r,ri	1/2	3	1/2	3	1	2.5	1	2	1	2.7
and	r,r	1/2	2	1/2	2	1	1.75	1	2	1	2.7
or	r,r	1/2	2	1/2	2	1	1.75	1	2	1	2.7
xor	r,r	1/2	2	1/2	2	1	1.75	1	2	1	2.7
inc	r	1/2	1.5	1/2	1.5	1	1	1	2	1	3
dec	r	1/2	1.5	1/2	1.5	1	1	1	2	1	3
neg	r	1/2	2	1/2	2	1	2	1	2	1	2.7
not	r	1/2	2	1/2	2	1	1.7	1	2	1	2.7
imul	r,ri	14	1/4	14	1/4	10	1	4	1	4	1/2
mul	r	14	1/10	14	1/10	11	1/2	5	1/2	6 ³	1/3
div	r	70 ²		70 ²		80 ²	1/34			39	1/39
adc	r,ri	7-8	1/6	7-8	1/6	10	1/4	2	0.75	1	2
sbb	r,ri	7-8	1/6	7-8	1/6	10	1/4	2	0.75	1	2
shl	r,i	4	1	4	1	1	1.7	1	1	1	2.7
sal	r,i	4	1	4	1	1	1.7	1	1	1	2.7
shr	r,i	4	1	4	1	1	1.7	1	1	1	2.7
sar	r,i	4	1	4	1	1	1.7	1	1	1	2.7
shl	r,c	6	1/6	6	1/6	2	1/2	1	1	1	2.7
sal	r,c	6	1/6	6	1/6	2	1/2	1	1	1	2.7
shr	r,c	6	1/6	6	1/6	2	1/2	1	1	1	2.7
sar	r,c	6	1/6	6	1/6	2	1/2	1	1	1	2.7
shld	r,r,i	12	1/12	8	1/4	8	1/7	2	1/2	2	1/2
shrd	r,r,i	14	1/14	8	1/4	8	1/7	2	1/2	2	1/2
shld	r,r,c	12	1/12	7	1/4	9	1/8	2	1/2	3	1/3
shrd	r,r,c	14	1/14	7	1/4	9	1/8	2	1/2	3	1/3
rol	r,i	4	1/4	4	1/4	1	1	1	1	1	2.7
ror	r,i	4	1/4	4	1/4	1	1	1	1	1	2.7
rol	r,c	6	1/6	6	1/6	2	1/2	1	1	1	2.7
ror	r,c	6	1/6	6	1/6	2	1/2	1	1	1	2.7
cmp	r,ri	1/2	3	1/2	3	1	2.5	1	2	1	2.7
test	r,i	1/2	2	1/2	2	1	1.7	1	2	1	2.7
test	r,r	1/2	2	1/2	2	1	1.7	1	2	1	2.7
bt	r,i	6	1/6	6	1/6	8	1/8	1	1	1	3
mov	r,r	1/2	3	1/2	3	1	2.5	1	2	1	2.7
cmov	r,r	6	1	6	1	10	1	2	1/2	1	2.3
movzx	r,r	1/2	3	1/2	3	1	2.5	1	2	1	2.3
movsx	r,r	1/2	2	1/2	2	1	1.75	1	2	1	2.3
bswap	r	7	0.67	7	0.67	1	2	2	1	1	2.7
lea	r,r+r	1/2	3	1/2	3	1	2.5	1	1	2	2.3
lea	r,r+r*s	3	1	3	1	1	1.25	1	1	2	2.3
lea	r,b+r+r	1	1.5	1	1.5	2	1.17	1	1	2	3
lea	r,b+r+r*s	4	1	4	1	2	0.8	1	1	2	3
bsr	r,r	8	1/2	8	1/2	16	1/2	2	1	9	1/8
bsf	r,r	8	1/2	8	1/2	16	1/2	2	1	7	1/7

Table 4. Integer register instructions, 32-bit operations. Please note that this table lacks more recent CPUs as well as some instructions which are present in the previous table.

		Intel P4		Intel P6		AMD	AMD	AMD	Intel
		F0-F2	F3-F4	D-E	F,17	K7	K8-K9	K10	Atom
		L T	L T	L T	L T	L T	L T	L T	L T
paddq	xmm, xmm	4 1/2	5 1/2	2 1/2	2 1	- -	2 1	2 2	5 1/5
psubq	xmm, xmm	4 1/2	5 1/2	2 1/2	2 1	- -	2 1	2 2	5 1/5
padd	xmm, xmm	2 1/2	2 1/2	1 1	1 2	- -	2 1	2 2	1 2
psubd	xmm, xmm	2 1/2	2 1/2	1 1	1 2	- -	2 1	2 2	1 2
pmuludq	xmm, xmm	6 1/2	7 1/2	4 1/2	3 1	- -	3 1/2	3 1	5 1/2
psllq	xmm, xmm	2 1/2	2 1/2	2 1/2	1 1	- -	2 1	3 2	5 1/5
psrlq	xmm, xmm	2 1/2	2 1/2	2 1/2	1 1	- -	2 1	3 2	5 1/5
psllq	xmm, i	2 1/2	2 1/2	2 1/2	1 1	- -	2 1	2 2	1 1
psrlq	xmm, i	2 1/2	2 1/2	2 1/2	1 1	- -	2 1	2 2	1 1
pslldq	xmm, i	4 1/2	4 1/2	3 1/3	2 1	- -	2 1	3 2	1 1
psrldq	xmm, i	4 1/2	4 1/2	3 1/3	2 1	- -	2 1	3 2	1 1
pand	xmm, xmm	2 1/2	2 1/2	1 1	1 3	- -	2 1	2 2	1 2
pandn	xmm, xmm	2 1/2	2 1/2	1 1	1 3	- -	2 1	2 2	1 2
por	xmm, xmm	2 1/2	2 1/2	1 1	1 3	- -	2 1	2 2	1 2
pxor	xmm, xmm	2 1/2	2 1/2	1 1	1 3	- -	2 1	2 2	1 2
movq	xmm, xmm	2 1/2	2 1/2	1 1	1 3	- -	2 1	2.5 3	1 2
punpckldq	xmm, xmm	2 1/2	2 1/2	2 1/2	2 1/2	- -	2 1	3 2	1 1
psadbw	xmm, xmm	4 1/2	4 1/2	4 1/2	3 1	- -	3 1/2	3 1	5 1/2
pshufd	xmm, xmm, i	4 1/2	4 1/2	2 1/2	4 1	- -	3 0.67	3 2	1 1

Table 5. Subset of integer SSE instructions.

		Intel P4		Intel P6		AMD	AMD	AMD	Intel
		F0-F2	F3-F4	B D	F	K7	K8-K9	K10	Atom
		L T	L T	L T	L T	L T	L T	L T	L T
paddq	mm, mm	2 1	2 1	- 2 - 1	2 1	- -	2 2	2 2	5 1/5
psubq	mm, mm	2 1	2 1	- 2 - 1	2 1	- -	2 2	2 2	5 1/5
padd	mm, mm	2 1	2 1	1 1	1 2	2 2	2 2	2 2	1 2
psubd	mm, mm	2 1	2 1	1 1	1 2	2 2	2 2	2 2	1 2
pmuludq	mm, mm	6 1	7 1	- 4 - 1	3 1	- -	3 1	3 1	4 1
pmul*w	mm, mm	6 1	7 1	- 3 - 1	3 1	- -	3 1	3 1	4 1
pmaddwd	mm, mm	6 1	7 1	3 1	3 1	- -	3 1	3 1	4 1
psllq	mm, mm	2 1	2 1	1 1	1 1	2 2	2 2	2 2	5 1/5
psrlq	mm, mm	2 1	2 1	1 1	1 1	2 2	2 2	2 2	5 1/5
psllq	mm, i	2 1	2 1	1 1	1 1	2 2	2 2	2 2	1 1
psrlq	mm, i	2 1	2 1	1 1	1 1	2 2	2 2	2 2	1 1
pand	mm, mm	2 1	2 1	1 1	1 3	2 2	2 2	2 2	1 2
pandn	mm, mm	2 1	2 1	1 1	1 3	2 2	2 2	2 2	1 2
por	mm, mm	2 1	2 1	1 1	1 3	2 2	2 2	2 2	1 2
pxor	mm, mm	2 1	2 1	1 1	1 3	2 2	2 2	2 2	1 2
movq	mm, mm	6 1	7 1	1 2	1 3	2 2	2 2	2 2	1 2
punpckldq	mm, mm	2 1	2 1	1 1	1 1	2 2	2 2	2 2	1 1
psadbw	mm, mm	4 1	4 1	5 4 1/2 1	3 1	3 1	3 1	3 1	4 1
pshufw	mm, mm, i	2 1	2 1	1 1	1 1	2 2	2 2	2 2	1 1

Table 6. Subset of integer MMX instructions.

Table remarks:

1. The latency is the indicated 5 cycles for the upper product half. The latency for the lower part is 4 cycles.
2. The latency is data dependent, the given numbers represent the worst case.
3. The latency is the indicated 6 cycles for the upper product half. The latency for the lower part is 4 cycles.
4. The latency is 77 cycles for model 0x17.
5. The latency is the indicated 10 cycles for the upper product half. The latency for the lower part is 3 cycles.
6. The latency is the indicated 4 cycles for the upper product half. The latency for the lower part is 3 cycles.
7. The latency is the indicated 2 cycles for the register result; the latency for the carry bit is just 1 cycle. The immediate operand 0 is handled specially; it gives a latency of just 1 also to the result register.